3. The classical oligopoly game

In this section we will discuss a special economic game with sets of strategies

$$X_{k} = [0, L_{k}] (L_{k} > 0, k=1,2,...,n)$$
 (19)

and pay-off functions

$$\Psi_{k}\left(x_{1},\ldots,x_{n}\right) = x_{k} f\left(\sum_{i=1}^{n} x_{i}\right) - K_{k}(x_{k}), \qquad (20)$$

where the functions f and K_k must have the properties: $\Re(f) = [0, L]$, where $L = \sum_{i=1}^{n} L_i$; $\Re(K_k) = [0, L_k]$;

 $R(f) \subset R^1$ and $R(K_k) \subset R^1$. The game defined by the sets of strategies (19) and pay-off functions (20) is called the <u>classical oligopoly game</u>.

Before discussing the equilibrium problem of this game we show how the game appears in some applications.

Application 1. Assume that n factories manufacture the same product and they sell it on the same market. Let f be the unit price of the product being a function of the total production level, and let K_k be the cost function of the manufacturer k. Then L_k is the production bound for manufacturer k and $\Psi_k(x_1, \dots, x_n)$ is its netto income assuming that x_i is the

production level of the manufacturer i for i=1,2,...,n .

<u>Application 2.</u> Assume that a multipurpose water supply system has to be designed. Let the water users denoted by k (k=1,2,...,n) and let the water quantity given to user k be denoted by x_k . If the capacity bounds of the users are denoted by L_k , then obviously $x_k \in [0, L_k]$ for $k=1,2,\ldots,n$. Let I be the investment cost being a function of $\sum_{k=1}^{n} x_k$, let $u_k(x_k)$, k=1

 $v_k(x_k)$ and $w_k(x_k)$ be the production cost, income and the economic loss of the water shortage /penalty e.t.c./ of user k, respectively. Let us assume, that the total investment cost is devided by the users in the rate of the water quantity used by the water users. Thus the total income of user k can be determined by the function

(n)

$$-\frac{\mathbf{x}_{k}}{\sum_{i=1}^{n} \mathbf{x}_{i}} I\left(\sum_{i=1}^{n} \mathbf{x}_{i}\right) - u_{k}(\mathbf{x}_{k}) + v_{k}(\mathbf{x}_{k}) - w_{k}(\mathbf{x}_{k}).$$
(21)
i=1

By introduceing the notations

$$f\left(\sum_{i=1}^{n} x_{i}\right) = -\frac{1}{\sum_{i=1}^{n} x_{i}} I\left(\sum_{i=1}^{n} x_{i}\right)$$
$$K_{k}(x_{k}) = u_{k}(x_{k}) - v_{k}(x_{k}) + w_{k}(x_{k})$$

function (21) has immediately form (20).

<u>Application 3.</u> Let us now assume that n factories are on the bank of a river and they send a certain quantity of

waste-water to the river. It is also assumed that the total

penalty paid by the factories is a function of the total waste-

-water quantity sent to the river and it is devided among the

factories proportionally to the waste-water quantity sent to the

river by the different factories. Let Lk be the total waste-

-water quantity produced by factory k, let x_k be the waste-water quantity sent to the river by factory k. Then the total "income" of factory k can be given by the formula

$$\frac{-\mathbf{x}_{k}}{\sum_{i=1}^{n} \mathbf{x}_{i}} P\left(\sum_{i=1}^{n} \mathbf{x}_{i}\right) - C_{k} (\mathbf{L}_{k} - \mathbf{x}_{k}), \qquad (22)$$

$$i=1$$

where Pis the penalty function, C_k is the cleaning cost of factory k. Let

$$f\left(\sum_{i=1}^{n} x_{i}\right) = -\frac{1}{\sum_{i=1}^{n} x_{i}} P\left(\sum_{i=1}^{n} x_{i}\right), \quad K_{k}(x_{k}) = C_{k}(L_{k} - x_{k}),$$
$$\lim_{i=1}^{n} x_{i} = 0$$

then the function (22) immidiately has the form of (20).

First we show that the equilibrium problem of the classical oligopoly game is equivalent to a fixed point problem of a one dimension point-to-set mapping. It will be much more convenient than the application of the fixed point problem of Lemma 3, since the latter is an n-dimensional problem.

Let

$$\begin{split} \Psi_k(\mathbf{s}, \mathbf{x}_k, \mathbf{t}_k) &= \mathbf{t}_k \ \mathbf{f} \left(\mathbf{s} - \mathbf{x}_k + \mathbf{t}_k \right) - \mathbf{K}_k(\mathbf{t}_k) \\ \text{for } k=1,2,\ldots,n, \ \mathbf{s} \in [0, \ L], \ \mathbf{x}_k \in [0, \ \mathbf{L}_k] \text{ and } \mathbf{t}_k \in [0, \mathbf{X}_k], \\ \text{where } \mathbf{X}_k &= \min \left\{ \mathbf{L}_k, \ \mathbf{L} - \mathbf{s} + \mathbf{x}_k \right\}. \text{ Since } \mathbf{X}_k &\cong 0, \text{ the interval} \\ \text{for } \mathbf{t}_k \text{ can not be empty. For } k=1,2,\ldots,n \ ; \ \mathbf{s} \in [0, \ L]; \ \mathbf{x}_k \in [0, \ \mathbf{L}_k] \end{split}$$





$$= \max_{0 \le u_k \le r_k} \Psi_k(s, x_k, u_k) \}$$

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and for k=1,2,...,n; s
$$\in [0, L]$$
 let
 $X_k(s) = \{x_k \mid x_k \in [0, L_k], x_k \in T_k(s, x_k)\}.$

Lemma 7. A vector
$$\underline{\mathbf{x}}^{\mathtt{H}} = (\mathbf{x}_{1}^{\mathtt{H}}, \dots, \mathbf{x}_{n}^{\mathtt{H}})$$
 is an equilibrium point of the classical oligopoly game if and only if $\mathbf{x}_{k}^{\mathtt{H}} \in \mathbf{X}_{k}(\mathbf{s}^{\mathtt{H}})$
(k=1,2,...,n), where $\mathbf{s}^{\mathtt{H}} = \sum_{k=1}^{n} \mathbf{x}_{k}^{\mathtt{H}}$.

<u>Proof.</u> The definition of the equilibrium point implies that a strategy vector $\underline{x}^{\underline{x}} = (x_1^{\underline{x}}, \dots, x_n^{\underline{x}})$ is an equilibrium point if and only if

$$x_{k}^{\mathtt{H}} f\left(s^{\mathtt{H}} - x_{k}^{\mathtt{H}} + x_{k}^{\mathtt{H}}\right) - K_{k}(x_{k}^{\mathtt{H}}) \geq t_{k} f\left(s^{\mathtt{H}} - x_{k}^{\mathtt{H}} + t_{k}\right) - K_{k}(t_{k})$$
(23)
for k=1,2,...,n and $t_{k} \in [0, L_{k}]$. /It is easy to observe that
for $s^{\mathtt{H}} = \sum_{i=1}^{n} x_{i}^{\mathtt{H}}$, $Y_{k} = L_{k}$ / Inequality (23) is equivalent
to the fact that $x_{k}^{\mathtt{H}} \in T_{k}(s^{\mathtt{H}}, x_{k}^{\mathtt{H}})$, that is $x_{k}^{\mathtt{H}} \in X_{k}(s^{\mathtt{H}})$.
Let us finally introduce the following one dimensional
point-to-set mapping:
 $X(s) = \left\{u \mid u = \sum_{i=1}^{n} x_{i}, x_{i} \in X_{i}(s)\right\}$ (s $\in [0, L]$). (24)

Lemma 7. and definition (24) imply the following important result.

Theorem 3. A vector
$$\underline{x}^{\texttt{H}} = (x_1^{\texttt{H}}, \dots, x_n^{\texttt{H}})$$
 is an equilibrium point of the classical oligopoly game if and only if for

$$s^{\#} = \sum_{i=1}^{n} x_{i}^{\#}$$
, $s^{\#} \in X(s^{\#})$ and for $k=1,2,\ldots,n$, $x_{k}^{\#} \in X_{k}(s^{\#})$.

Remark. The solution of the game has two steps: Step 1: the solution of the one dimensional fixed point problem $s^{\#} \in X(s^{\#})$; Step 2: the determination of sets $X_k(s^{\#})$ and the computation of the vectors $\underline{x}^{\#} = (x_1^{\#}, \dots, x_n^{\#})$ such that $x_k^{\#} \in X_k(s^{\#})$ $(k=1, 2, \dots, n)$ and $s^{\#} = \sum_{k=1}^n x_k^{\#}$.

In the following parts of this section we will assume that the conditions given below are satisfied.

1. There exists a constant $\xi > 0$ such that

a/ f(s) = 0 for $s \ge \xi$; b/ f is continuous, concave and strictly decreasing in the interval $[0, \xi]$.

2. For $k=1,2,\ldots,n$ function K_k is continuous, convex and strictly increasing in the interval $[0, L_k]$.

Theorem 4. Under the above conditions the game has at least one equilibrium point.

<u>Proof.</u> The proof consists of several steps. a/ First we prove that if $\underline{x}^{\texttt{H}} = (x_1^{\texttt{H}}, \dots, x_n^{\texttt{H}})$ is an equilibrium point, then $\sum_{k=1}^{n} x_k^{\texttt{H}} \leq \xi$. Let us suppose that





$$\begin{aligned} \Psi_k \left(x_1^{\texttt{H}}, \dots, x_k^{\texttt{H}}, \dots, x_n^{\texttt{H}} \right) &= x_k \cdot 0 - K_k(x_k) > x_k^{\texttt{H}} \cdot 0 - K_k(x_k^{\texttt{H}}) = \\ \Psi_k \left(x_1^{\texttt{H}}, \dots, x_k^{\texttt{H}}, \dots, x_n^{\texttt{H}} \right), & \text{which is a contradiction to inequality} \\ (1). \end{aligned}$$

b/ Let $X = \left\{ \underline{x} \mid \underline{x} = (x_1, \dots, x_n), \sum_{k=1}^n x_k \leq \xi, x_k \notin [0, L_k], \\ k=1,2,\dots,n \right\}.$ Next we prove that any equilibrium point $\underline{x}^{\underline{\pi}}$ of the generalized game $\Gamma = (n; X_1, \dots, X_n, X; \mathcal{C}_1, \dots, \mathcal{C}_n)$ gives an equilibrium point for the classical oligopoly game. Let $x_k \notin [0, L_k].$ If $(x_1^{\underline{\pi}}, \dots, x_{k-1}^{\underline{\pi}}, x_k, x_{k+1}^{\underline{\pi}}, \dots, x_n^{\underline{\pi}}) \notin X$, then the equilibrium property for game Γ gives

$$\begin{aligned} \Psi_k \left(x_1^{\texttt{H}}, \dots, x_k^{\texttt{H}}, \dots, x_n^{\texttt{H}} \right) &\geq \Psi_k \left(x_1^{\texttt{H}}, \dots, x_k^{\texttt{H}}, \dots, x_n^{\texttt{H}} \right), \\ \text{and if } \left(x_1^{\texttt{H}}, \dots, x_k^{\texttt{H}}, \dots, x_n^{\texttt{H}} \right) \not\in \mathbb{X}, \text{ then} \end{aligned}$$

$$\Psi_k(x_1^{H}, \dots, x_k^{H}, \dots, x_n^{H}) = x_k \cdot 0 - K_k(x_k) < - K_k(0) =$$

$$= 0.f\left(\sum_{i \neq k}^{\mathcal{H}} x_{i}^{\mathcal{H}}\right) - K_{k}(0) = \mathcal{I}_{k}\left(x_{1}^{\mathcal{H}}, \dots, 0, \dots, x_{n}^{\mathcal{H}}\right) \leq \mathcal{I}_{k}\left(x_{1}^{\mathcal{H}}, \dots, x_{k'}^{\mathcal{H}}, x_{n}^{\mathcal{H}}\right)$$

since $\left(x_{1}^{\mathcal{H}}, \dots, 0, \dots, x_{n}^{\mathcal{H}}\right) \notin X.$

c/ Next we prove that if function h is continuous, concave and strictly decreasing in a nonnegative interval [A, B], then the function xh(x) is concave in the same interval.

Let us first assume that h is twice continuously

differentiable.

Then

$$\{xh(x)\}' = xh'(x) + h(x),$$

$$\{xh(x)\}'' = 2h'(x) + xh''(x) < 0,$$

which implies the assertion.

If h is continuous, then let h_m (m=1,2,...) be twice continuously differenciable, concave, strictly decreasing functions such that $\lim h_m = h_o$ m > 00

Let $A \leq x < y \leq B$; $x, \beta \geq 0$; $x - \beta = 1$, then for m = 1, 2...

$$(\alpha x + \beta y)h_{m} (\alpha x + \beta y) \geq \alpha xh_{m}(x) + \beta yh_{m}(x).$$

By the limit relation $m \rightarrow \infty$ we obtain

 $(\alpha x + \beta y)h(\alpha x + \beta y) \ge \alpha xh(x) + \beta yh(y),$

thus xh(x) is concave.

d/ The parts a/ and b/ imply that the classical oligopoly game and the generalized game $\Gamma = (n; X_1, \dots, X_n, X; Y_1, \dots, Y_n)$ have the same equilibrium points. Under the assumptions of the theorem X is a convex, closed, bounded subset of \mathbb{R}^n , Ψ_k is continuous and part c/ implies that Ψ_k is concave in x_k . Thus the conditions of the Nikaido-Isoda theorem are satisfied, consequently the game has at least one equilibrium point.

not Remark. The uniqueness of the equilibrium point is Vassured in general as the following example shows.



$$f(s) = \begin{cases} 1,75 - 0,5s , & \text{if } 0 \le s \le 1,5 \\ 2,5 - s , & \text{if } 1,5 \le s \le 2,5 \\ 0 & , & \text{if } s > 2,5 ; \end{cases}$$

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$$K_1(x) = K_2(x) = 0,5x \quad (x \ge 0).$$

We will prove that an arbitrary point of the set

$$\mathbf{X}^{\mathbf{x}} = \left\{ (\mathbf{x}_1, \mathbf{x}_2) \middle| 0, 5 \leq \mathbf{x}_1 \leq 1, 0, 5 \leq \mathbf{x}_2 \leq 1, \mathbf{x}_1 + \mathbf{x}_2 = 1, 5 \right\}$$

gives an equilibrium point of the game.

Lét x [0,5;1] be fixed, and let

$$\Psi(x) = xf(1,5 - x + x^{m}) - K_{k}(x) \quad (k=1,2,).$$

It is easy to verify that

$$\Psi'(x^{\underline{\pi}} - 0) = x^{\underline{\pi}}(-0,5) + 1 - 0,5 = 0,5(1 - x^{\underline{\pi}}) \ge 0$$
,

and

$$\Psi'(x^{\mathbf{X}}+0) = x^{\mathbf{X}}(-1)+1 - 0,5 = 0,5 - x^{\mathbf{X}} \leq 0.$$

Part c/ implies that function Ψ is concave in x, consequently from the inequalities $\Psi'(x^{\pm} - 0) \geq 0$ and $\Psi'(x^{\pm} + 0)$ we can conclude that x^{\pm} is a maximum point of the function Ψ . Thus arbitrary $x^{\pm} \in X^{\pm}$ is an equilibrium point.

Next we discuss a numerical algorithm for findig the equilibrium points of the classical oligopoly game. Under the assumptions of Theorem 4. the following statements are true.

Lemma 8.

a/ For s $\in [0, L]$, $X_{i}(s)$ is not empty and is a closed interval $\left[A_k(s), B_k(s)\right]$, $\left[k=1, 2, \ldots, n\right]$;

b/ for $0 \le s < s' \le L$ the inequality $B_k(s') \le A_k(s)$ holds for k=1,2,...,n;

c/ if f is differentiable at the point s, then $A_k(s) = B_k(s)$;

d/ if f is differentiable in the interval [0, L], then $A_k(s)$ is a continuous function of s.

<u>Proof.</u> Parts a/ and b/ can be proven by simple modifications of parts C/a/ and C/b/ of the proof of Theorem 1. in paper [10]. The statements c/ and d/ are proven in the C/a,b,c

part of the proof of Theorem 1. in paper [10].

Lemma 9. If $\underline{x}^{\texttt{M}} = (\underline{x}_{1}^{\texttt{M}}, \dots, \underline{x}_{n}^{\texttt{M}})$ and $\underline{x}^{\texttt{M}} = (\underline{x}_{1}^{\texttt{M}}, \dots, \underline{x}_{n}^{\texttt{M}})$ are equilibrium points of the classical oligopoly game having the properties given in Theorem 4., then



which is a contradiction.

Corollary. The point-to-set mapping X(s) has exactly one

fixed point, which can be computed by the usual bisection method (see F. Szidarovszky, S.Yakowitz [12]).

<u>Theorem 5</u>. Assume that the conditions of Theorem 4. are satisfied. Let $s^{\mathbb{H}}$ be the unique fixed point of the mapping X(s). Then all equilibrium points of the classical oligopoly game can be obtained by the solution of the system of linear equations and inequalities:

$$A_{k}(s^{\underline{w}}) \leq x_{k} \leq B_{k}(s^{\underline{w}}) \quad (k=1,2,\ldots,n)$$
$$\sum_{k=1}^{n} x_{k} = s^{\underline{w}}$$
$$k=1$$

Proof. The statement is a consequence of Lemma 8. and

Lemma 9.



<u>Corollary</u>. If in addition to the conditions of Theorem 4. function f is differentiable on the interval [0, L], then the equilibrium point is unique.

Remark 1. It is interesting to observe that the game is not linear but the set of equilibrium points is a simplex.

Remark 2. The uniqueness of the equilibrium point depends on the differentiability of a function and not on strict concavity as it is usual in the theory of nonlinear programming.

Special cases.

1. In case of f and K_k $(l \leq k \leq n)$ being twice differentiable the uniqueness was proved by 0.0pitz [7] without giving any

algorithm for finding it.

2. Under the assumptions of 0. Opitz, F. Szidarovszky [9] proved the existence and uniqueness of the equilibrium point

and also gave an iterative algorithm for computing it.

3. If the cost functions K_k are identical and the conditions of 0.0pitz are satisfied, then **E**. Burger [1] proved the existence and uniqueness of the equilibrium point and also gave an algorithm to compute it. We remark that the algorithm of Szidarovszky is a generalization of Burger's method.

4. If the functions f and K_k (k=1,2,...,n) are linear, then the existence and uniqueness was proved by M. Maňas,[4], who gave an algorithm which is independent of the method of Szidarovszky. We remark that using the result of Theorem 5.

the equilibrium point in this special case can be given in closed form (see pp. 37-39 of [10]).

4. The group equilibrium problem

In this paragraph we will discuss the generalized version of the classical oligopoly game Γ having the stategy sets $X_{k} = [0, L_{kl}] \times [0, L_{k2}] \times \cdots \times [0, L_{ki_{k}}]$ ($l \le k \le n$) (25) and pay-off functions $\Psi_{k}(\underline{x}_{l}, \ldots, \underline{x}_{n}) = \left(\sum_{i=1}^{i_{k}} x_{ki}\right) f\left(\sum_{\ell=1}^{n} \sum_{j=1}^{i_{\ell}} x_{\ell j}\right) - K_{k}(\underline{x}_{k}),$ (26)

where for $k=1,2,\ldots,n$, $\underline{x}_{k} = (x_{k1},\ldots,x_{ki_{k}}) \in \underline{x}_{k}$. This game

can occur when the players of the classical oligopoly game form disjoint groups and they tend to the optimal income of the group. If the number of members in group k is equal to i_k ,